# HOW TO FACTOR OBJECTS?

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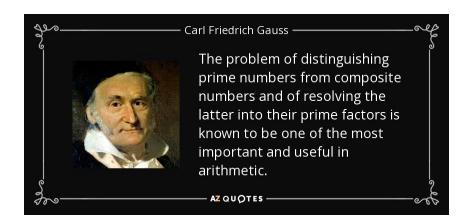
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## BASE CASES

### INTEGERS

- Integer n factors uniquely into prime numbers.
  - $\triangleright$  Eg. 1092 =  $2^2 \times 3 \times 7 \times 13$

- ❖ Given n, can you factor it?
  - Input n in binary
  - ➤ 2<sup>^</sup>{(log n)<sup>0.3</sup>} time not good enough
  - Number Field Sieve (1990s) factors via x² = y² mod n
- Hardness used in cryptosystems.
  - RSA, HTTPS, SSh, SFTP, Diffie-Hellman, Banks, Whatsapp, ...



## Prime Numbers Eratosthenes'(ehr-uh-TAHS-thuh-neez) Sieve



- •Eratosthenes was a Greek mathematician, astronomer, geographer, and librarian at Alexandria, Egypt in 200 B.C. •He invented a method for finding prime numbers that is still used today.
- This method is called Eratosthenes' Sieve.

276 BC - 194 BC

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### UNIVARIATE OVER INTEGERS

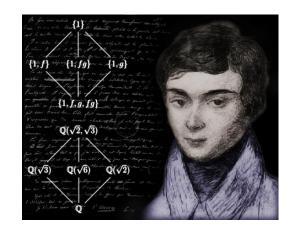
- $\diamond$  Given polynomial  $f(x) \in \mathbb{Z}[x]$ , factor it.
  - $\rightarrow$  f =  $x^5-x^4-4x^2+x-2$  factors
  - > Roots have no formula
  - Irreducibility testing?



- [Lenstra, Lenstra, Lovász'82] solved this completely.
  - ightharpoonup Factor mod 2, 2<sup>2</sup>, 2<sup>4</sup>, 2<sup>8</sup>,...
  - ➤ Lift to integral factor using lattice theory
  - Useful in many post-quantum cryptosystems

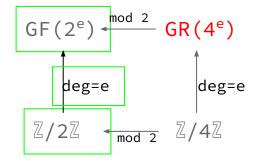
#### UNIVARIATE OVER FINITE FIELDS

- Galois field GF(pe) of size pe and char =
  prime p.
- $\bullet$  Given polynomial  $f(x) \in GF(p)[x]$ , factor it.
  - $\rightarrow$  f =  $x^2-2$  factors mod 7
  - $> \sqrt{2} = 3 \mod 7$ !
  - > Irreducibility testing?
- [Berlekamp'67; Cantor-Zassenhaus'81] solved this practically.
  - Use Galois automorphism
  - $\triangleright$  Compute gcd of f(x) with  $x^{p}-x$ ,  $x^{p^{2}}-x$ ,  $x^{p^{3}}-x$ ,...
  - Useful in crypto, coding theory, computational algebra, arithmetic-geometry, ...



## UNIVARIATE OVER GALOIS RINGS

- Galois ring GR(p<sup>ke</sup>) of size p<sup>ke</sup> and characteristic = prime-power p<sup>k</sup>.
- ❖ Given polynomial  $f(x) ∈ GR(p^{ke})[x]$ , factor it.
  - $\rightarrow$  f =  $x^2-2$  factors mod  $7^2$
  - $> \sqrt{2} = 10 \mod 7^2$ !
  - > Irreducibility testing?



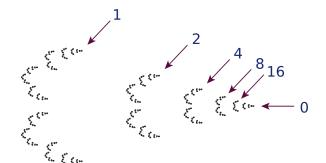
- This problem is OPEN.
- ♦ [Dwivedi, Mittal, S.'19] solved for k=4.
  - $\triangleright$  Factor  $f(x) \mod p$ ,  $p^2$ ,  $p^3$ ,  $p^4$ .
  - Lifting from one to the next precision is nontrivial.
  - $\triangleright$  Eg. f =  $x^2$ -p mod p<sup>2</sup> vs f =  $x^2$ -px mod p<sup>2</sup>

## UNIVARIATE OVER P-ADIC NUMBERS

- - > 1+2p+3p<sup>2</sup>+4p<sup>3</sup>+5p<sup>4</sup>+... converges to a number!
- ❖ Given polynomial  $f(x) ∈ Z_p[x]$ , factor it.
  - $\rightarrow$  f =  $x^2$ -2 factors in 7-adic
  - $> \sqrt{2} = 3 + 1*7 + 2*7^2 + 6*7^3 + \dots$  in infinite digits!
  - Irreducibility testing?



- [Chistov'90; Cantor, Gordon'00] solved it efficiently.
  - $\triangleright$  Newton polytope of f(x),
  - ➤ coupled with p-adic metric,
  - ➤ reduces to mod p factoring.
  - > Useful in computational number theory.

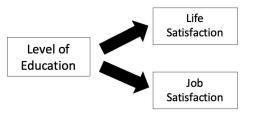


## MULTIVARIATES

## MULTIVARIATE SPARSE POLYNOMIALS

- Given polynomial  $f(x_1, x_2, ..., x_n) \in F[x]$ , factor it.
  - $\rightarrow$  f =  $(x_1^d-1)...(x_n^d-1)$  factors into
  - $\Rightarrow$  g =  $(x_1^{d-1} + ... + x_1 + 1) ... (x_n^{d-1} + ... + x_n + 1)$ .
  - $\triangleright$  Sparsity s:=2<sup>n</sup> blows-up to d<sup>n</sup>.
  - > => Factors can be very large!
- ❖ What if individual-degree d is constant?
- [Bisht,S.'22] showed a poly bound for symmetric factors.
  - $\triangleright$  Newton polytope of f(x)
  - > Relation between #vertices & #internal points.
  - Fast algorithm, by reducing to the base cases

#### Multivariate



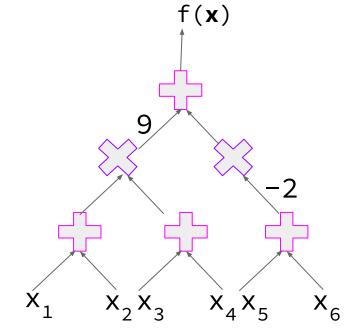
#### Univariate



## IN FORMULA MODEL

- ♦ Given polynomial  $f(x_1, x_2, ..., x_n) ∈ F[x]$ , factor it.
  - > Input: is a formula of size s.
  - > Output: is a formula of size =?
- Only quasipoly bound known till 2024.
- ♦ [BKRRSS'25] prove Poly bound in 2025.

Idea: Use an expression for the roots given by Lagrange Inversion formula of analytic functions.



Lagrange Inversion Theorem.

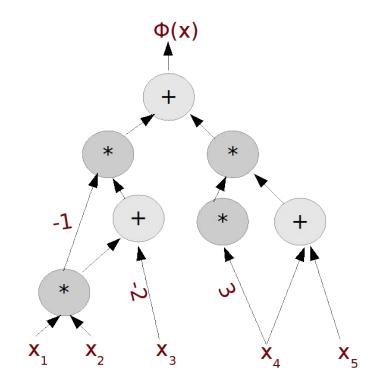
If a GF 
$$g(z) = \sum_{n \ge 1} g_n z^n$$
 satisfies the equation  $z = f(g(z))$  with  $f(0) = 0$  and  $f'(0) \ne 0$  then  $g_n = \frac{1}{n} [u^{n-1}] \left(\frac{u}{f(u)}\right)^n$ .

## IN CIRCUIT MODEL

- ❖ Given polynomial  $f(x_1, x_2, ..., x_n) ∈ F[x]$ , factor it.
  - ➤ Input: is a circuit of size s.
  - ➤ Output: is a circuit of size =?

- [Kaltofen'87] showed a poly bound.
  - degree not too `high'

- ❖ Idea: Same as formula.
  - ➤ Works with constant-depth circuits!

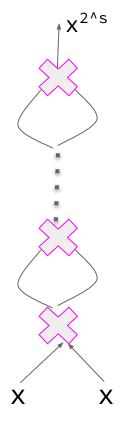


## IN A 'TOUGHER' CIRCUIT MODEL

- ♦ Given polynomial  $f(x_1, x_2, ..., x_n) \in F[x]$ , factor it.
  - ➤ Input: is a circuit of size s and degree 2<sup>s</sup>.
  - > Output: a factor of degree poly(s) of size =?

- It's an open question.
- \* [Dutta,S.,Sinhababu'18] showed a poly bound, when
  - degree of the radical of f is not too `high'.

Idea: all-roots-Newton-iteration is doable inside circuits.



## CONCLUDE WITH OPEN PROBLEMS

- Question 1: Fast integer factoring?
- ❖ Question 2: Fast polynomial factoring mod p<sup>k</sup>?

- Question 3: General circuit factoring?
- Question 4: Derandomization?



