

$\Rightarrow \Sigma := \bigcup_{\text{rk}(T) > \tau} \Sigma_T$  has size  $< p \cdot \epsilon^{-\tau/8q} \cdot q^n$

&  $A := \mathbb{F}_q^n \setminus \Sigma$  zeroes out every function in  $\partial^k T$ , for  $T \in \{T_i \mid i \in [s], \text{rk}(T_i) > \tau\}$ .

$\Rightarrow \Gamma_{k,A}^1(c)$  is contributed by only  $T_i$ 's with  $\text{rk}(T_i) < \tau$

$\Rightarrow \Gamma_{k,A}^1(c) < p \cdot q^\tau. \quad \square$

- Next, we understand the measure for  $\text{det}_d$  &  $\text{per}_d$ ,  $n := d^2$ .

Lemma 2 (Lower bound): For any  $A \subseteq \mathbb{F}_q^n$  of size  $(1 - o(1))q^n$ , we have  $\Gamma_{k,A}^1(\text{det}_d) = \binom{d}{k}^2$ .

Proof: (from Saptharishi's survey)

- We consider the rank of the matrix  $M_k(\text{det}_d, A)$ .

- An order- $k$  derivative (partial) of  $\text{det}_d$  is, either zero, or an order- $(d-k)$  minor.

• Since  $\det_d$  has  $\binom{d}{d-k}$  many order- $(d-k)$  minors, it can be seen that the rank of  $M_k(\det_d, \mathbb{F}_q^n) = \binom{d}{k}$ .

[ We can pick a point  $\bar{x} \in \mathbb{F}_q^n$  s.t. the column  $\bar{x}$  has exactly one nonzero entry in  $M_k(\det_d, \mathbb{F}_q^n)$ : a desired order- $(d-k)$  minor. Thus, we identify a "diagonal" matrix inside  $M_k(\cdot, \cdot)$ ; lower bounding its rank. ]

• However,  $M_k(\det_d, A)$  has, possibly, many columns missing. How do we lower bound its rank?

Idea - We study arbitrary linear combinations of its rows.

Claim 1: Let  $f(\bar{x})$  be a  $\mathbb{F}_q$ -linear combination of  $r \times r$  minors of  $X = (x_{ij})$ . Then,

$$\Pr_{\bar{x} \in \mathbb{F}_q^n} [f(\bar{x}) \neq 0] \geq \frac{1}{4}.$$

• This claim immediately implies that the rows of  $M_k(\det_d, A)$ , corresponding to the minors, are linearly independent; since the #zeros, of a linear combination of minors of  $X$ , is  $\leq \frac{3}{4}q^n$  & so  $|A| - \frac{3}{4}q^n = \frac{1}{4}q^n - o(1)q^n > 0$ .

$\Rightarrow$  We only need to prove Claim 1. First, we prove a base case:

Claim 2:  $\Pr_{\bar{x} \in \mathbb{F}_q^n} [\det_d(\bar{x}) \neq 0] \geq 1/4$ .

Proof: • The number of invertible matrices in  $\mathbb{F}_q^{d \times d}$  is  $(q^d - 1) \cdot (q^d - q) \cdot \dots \cdot (q^d - q^{d-1})$ .

$$\begin{aligned}
 \Rightarrow \Pr_{\bar{x}} [\det_d(\bar{x}) \neq 0] &= \left(1 - \frac{1}{q}\right) \cdot \left(1 - \frac{1}{q^2}\right) \cdot \dots \cdot \left(1 - \frac{1}{q^d}\right) \\
 &\geq \left(1 - \frac{1}{2}\right) \left(1 - \frac{1}{2^2}\right) \cdot \dots \cdot \left(1 - \frac{1}{2^d}\right) \geq \frac{1}{4}.
 \end{aligned}$$

Exercise: Prove Claim 2 for  $\text{per}_d$ .

Pf of Claim 1: • Let the linear combination of the  $r \times r$  minors of  $\det_d(X)$  be

$$f(\bar{x}) = \sum_{\text{row-1 in } M_i} c_i \cdot M_i + \sum_{\text{row-1 not in } M_j} c_j \cdot M_j.$$

• We now want to further expand each  $M_i$  by row-1 of  $X$  & rearrange the first part of  $f(\bar{x})$  above:

$$f(\bar{x}) = \sum_{i \in [d]} x_{1i} \cdot M'_i + M''.$$

Now  $M'_i$  are  $\mathbb{F}_q$ -linear combinations of certain order- $(r+1)$  minors of  $\det_d(X)$ .

$M''$  is "free" of  $x_{1j}$  variables.

• Wlog we can assume that at least two distinct order- $r$  minors participated in defining  $f(\bar{x})$ , and that at least one of the  $M'_i$  above is nonzero.

• We would like to pick a random  $\bar{x}$  by first picking the rows  $\{2, \dots, d\}$  & picking the

- first row in the end (from  $\mathbb{F}_2^d$ ).
- From this viewpoint it is clear that:

$$\text{LHS} = \Pr_{\alpha} \left[ \sum_{i=1}^d \alpha_{1i} \cdot M_i' + M'' \neq 0 \right]$$

$$\geq \Pr_{\alpha} \left[ \sum_{i=1}^d \alpha_{1i} \cdot M_i' \neq 0 \right] \quad (\text{koufis' trick}),$$

- The latter involves only the minors that have row 1 of  $X$ .
- Repeating this, several times, we end up with the probability estimate for a single minor (as in Clm 2).  
 $\Rightarrow \text{LHS} \geq 1/4$ . □

- As discussed before Clm 1 implies that  $\prod_{k,A} (\det_d) = \binom{d}{k}^2$ , finishing Lem 2. □

Exercise: Prove the same for  $\text{per}_d$ .

• Assuming that  $\det_d$  has a depth-3 circuit, we compare the bounds in Lemmas 1 & 2: let  $\tau = \alpha d$ ,  $k = \tau / 10q$ ,  

$$\binom{d}{k}^2 = \prod_{k, A} (\det_d) < \delta \cdot q^{\alpha d}$$

[ Stirling's approx. gives:  $\lg \binom{n}{\epsilon n} = H_2(\epsilon) \cdot n - O(\lg n)$ ,

where  $H_2(\epsilon) := -\epsilon \lg \epsilon - (1-\epsilon) \lg (1-\epsilon)$ . ]

[ e.g. it follows that  $\binom{n}{\epsilon n} = 2^{\Omega_\epsilon(n)}$ . ]

$$\Rightarrow \lg \binom{d}{k}^2 = \Omega(d \cdot H_2(k/d)) = \Omega(d \cdot H_2(\alpha/10q))$$

$$\Rightarrow \lg s = \Omega\left(d H_2\left(\frac{\alpha}{10q}\right)\right) - \alpha d \lg q$$

$$\Rightarrow \lg s / d = \Omega\left(\frac{\alpha}{10q} \lg \frac{10q}{\alpha} + \left(1 - \frac{\alpha}{10q}\right) \lg \frac{10q}{10q - \alpha}\right) - \alpha \cdot \lg q$$

• Thus, there is some constant  $c > 0$  s.t. it suffices to pick  $\alpha$  satisfying

$$\lg \frac{10q}{\alpha} > cq \cdot \lg q$$

$$\Leftrightarrow \alpha < 10q / q^{c-1}. \text{ Thus, } \tau = O(d / q^{c-1}).$$

- For constant  $q$ ,  $\tau$  makes sense & we get a lower bound on the top fanin:  

$$\lg s = \Omega_q(d)$$
 finishing the theorem.  $\square$

- The lower bound can be improved by considering a sum of elementary symmetric polynomials on  $n = d^2$  variables &  $\deg \leq d$ .

$$\text{Define } \underline{\text{sym}}_{\leq d} := \sum_{SE\left(\begin{matrix} [n] \\ \leq d \end{matrix}\right)} x_s.$$

- It can be shown that the rank of the matrix  $M_k(\text{sym}_{\leq d}, \mathbb{F}_q^n) \geq \binom{n}{d/2}$ , for  $k = d/2$ .

- This gives  $s = n^{\Omega_q(d)}$ .