## CS747 - RANDOMIZED METHODS IN COMPUTATIONAL COMPLEXITY NITIN SAXENA

## **ASSIGNMENT 2**

POINTS: 50

DATE GIVEN: 05-SEP-2025 DUE: 25-SEP-2025

## Rules:

• You are strongly encouraged to work *independently*. That is the best way to understand & master the subject.

- Write the solutions on your own and honorably *acknowledge* the sources if any. http://cse.iitk.ac.in/pages/AntiCheatingPolicy.html
- Submit your solutions, before time, to your TA. Please give your LaTeXed or Word processed solution-sheet in PDF. This will be graded, and commented, in-place.
- Clearly express the fundamental *idea* of your proof/ algorithm before going into the other proof details. The distribution of partial marks is according to the proof steps.
- There will be a penalty if you write unnecessary or unrelated details in your solution. Also, do not repeat the proofs done in the class.
- Problems marked '0 points' are for practice.
- Acknowledgements: Several problems are from *Arora & Barak, Computational Complexity: A Modern Approach*, and other lecture notes.

Question 1: [6+3 points] In the class we proved the Razborov-Smolensky result:  $\text{mod}_p \notin ACC^0[q]$ , for p=2=q-1. Give the general proof for primes  $p \neq q$ .

What happens to the proof when one of them is a composite?

**Question 2:** [7 points] Show that the OR of n variables cannot be expressed as a polynomial over  $\mathbb{F}_p$  of degree less than n.

Question 3: [13 points] Give an example of a set-family  $\mathcal{Z}$  with  $(p-1)^{\ell}$  subsets (each size  $\leq \ell$ ), such that there is no sunflower with p-petals.

**Question 4:**  $[6+9 \ points]$  Let A be a real symmetric stochastic matrix. Show that

- (1) its eigenvalues are real with absolute value at most 1.
- (2) its eigenvectors, for different eigenvalues, are linearly independent.

Question 5: [6 points] Let A, B be real symmetric stochastic matrices. Show that the second-largest eigenvalue function  $\lambda(\cdot)$  is subadditive (i.e.  $\lambda(A+B) \leq \lambda(A) + \lambda(B)$ ).

**Question 6:** [0 points] Let G be a graph with only two distinct eigenvalues. Prove that G is a complete graph.

Question 7: [0 points] Show that for an  $n \times n$  matrix  $X = ((x_{i,j}))$ , the polynomial per(X) has a formula of size  $2^{O(n)}$ .

**Question 8:** [0 points] Show that AM[k] = AM[2] for all constant  $k \geq 2$ .

**Question 9:** [0 points] What big change happens in the above proof, about AM[k], if k is non-constant?

**Question 10:** [0 points] Show that the graph isomorphism problem is in NP $\cap$ coAM. (So, "almost" in NP $\cap$ coNP.)

Is there a better (algorithmic) result known for graph isomorphism?

**Question 11:** [0 points] Make a list of interesting problems that are in AM $\cap$ coAM. (So, "almost" in NP $\cap$ coNP.)

Could any such problem be NP-hard?

Question 12: [0 points] Fix parameters n, d. Show that a random n-vertex d-regular graph is a (n, 2d, 1/3)-combinatorial expander.

Question 13: [0 points] Consider a subset A of the n-tuples  $\mathbb{F}_3^n$ . We are interested in the largest set A that is free of arithmetic progressions of size 3 (i.e. there is no  $a, b, c \in A$  satisfying  $a + b + c \equiv \mathbf{0} \mod 3$ ). Could you relate this question to the sunflower lemma?

Question 14: [0 points] Recall our proof of the  $n^{\Omega(\sqrt{k})}$  monotone circuit lower bound for the k-clique problem. Is this bound tight?

 ${\it Question~15:~[0~points]}$  Think of another "hard" problem that is monotone.