CS747 - RANDOMIZED METHODS IN COMPUTATIONAL COMPLEXITY
NITIN SAXENA

ASSIGNMENT 4

POINTS: 50

DATE GIVEN: 01-NOV-2018 DUE: 15-NOV-2018

Rules:

- You are strongly encouraged to work *independently*. That is the best way to understand the subject.
- Write the solutions on your own and honorably *acknowledge* the sources if any. `http://cse.iitk.ac.in/pages/AntiCheatingPolicy.html`
- Submit your solutions, before time, to your TA. Preferably, give the TA a printed copy of your LaTeXed or Word processed solution sheet.
- Clearly express the fundamental *idea* of your proof/ algorithm before going into the other proof details. The distribution of partial marks is according to the proof steps.
- There will be a penalty if you write unnecessary or unrelated details in your solution. Also, do not repeat the proofs done in the class.
- Problems marked '0 points' are for practice.
- Acknowledgements: Several problems are from *Arora & Barak, Computational Complexity: A Modern Approach* and other lecture notes.

Question 1: *[9 points]* Suppose boolean function f is in E with $\mathrm{H}_{\mathrm{avg}}(f) \geq n^4$. Then, the function $g : z \mapsto z \circ f(z)$, for $z \in \{0,1\}^\ell$, is an $(\ell+1)$-prg.

Question 2: *[13 points]* Prove that NEXP = MA implies NEXP $\subseteq$ P/poly.

Question 3: *[10 points]* Suppose BPP$\neq$EXP. Could you use this to derandomize BPP to some extent? Sketch the proof details.

Question 4: *[9 points]* For every $\delta > 0$ and sufficiently large n, prove the existence of a *linear* ecc $E : \{0,1\}^n \to \{0,1\}^{1.1n/(1-H(\delta))}$ with distance at least δ.

(Note: *Linear code* means that $E(x) + E(y) = E(x+y)$, where addition is componentwise modulo 2. Let $H(\delta) := -\delta \log \delta - (1 - \delta) \log(1 - \delta)$.)

Question 5: *[9 points]* Show that there exists an *explicit linear-stretch* code, i.e. $\exists c, \delta > 0$, $\forall n$, $\exists$ ecc $E : \{0,1\}^n \to \{0,1\}^{cn}$ of distance at least δ with efficient encoding/decoding algorithms.

Question 6: *[0 points]* For every $c > 0$, prove that EXP$\not\subseteq$ i.o.-Size(n^c).

Question 7: *[0 points]* Prove that, for unique decoding, the channel error should be less than 25%.

Question 8: *[0 points]* How do you factor $f(x)$ mod p?

Question 9: *[0 points]* How do you factor $f(x_1, x_2)$ mod p?

Question 10: *[0 points]* How do you find an integral root of an integral polynomial $f(x)$?

□□□