

- Thus, we've shown for any smooth proj. curve C over \mathbb{F}_q , genus g , the roots α of $L(t)$ satisfy $|\alpha| = \sqrt{q}$.

- $-2g\sqrt{q} \leq N_1(C) - (q+1) \leq 2g\sqrt{q}$.

\Rightarrow For $q \gg g^2$, $N_1(C) \approx q+1 = N_1(\mathbb{P}^1)$.

- RTH has tons of implications. For eg. in CS we use:

Corollary (Weil estimate for χ -sums): Let $\chi = \chi_2: \mathbb{F}_q \rightarrow \{-1, 0, +1\}$ be the character

$\alpha \mapsto \alpha^{(q-1)/2} \equiv 1$ iff α is square in \mathbb{F}_q .

Let $f(x)$ be deg- s polynomial. Then,

$$\left| \sum_{\alpha \in \mathbb{F}_q} \chi(f(\alpha)) \right| \leq O_s(\sqrt{q}).$$

Pf: • Consider the curve C for $k := \mathbb{F}_q(x)[y]/\langle y^2 - f \rangle$.

$$\bullet \sum_{\alpha \in \mathbb{F}_q} \chi(f(\alpha)) = \sum_{\alpha: \chi(f(\alpha))=1} 1 - \sum_{\alpha: \chi(f(\alpha))=-1} 1$$

$$= \sum_{\alpha: \chi(f(\alpha))=1} 2 - \sum_{\alpha: \chi(f(\alpha))=-1, 1} 1 = N_1(C) - q + O_s(1) = O_s(\sqrt{q}). \quad \square$$

Exercise: Do this for other exponential sums.

Qn: (i) Given C/\mathbb{F}_q , how do we compute $N_1(C)$ in $\text{poly}(\log q)$ -time?

(ii) Is there another interpretation of $L(t)$ that can help in computing?

Cohomological interpretation of $L(t)$

- See Frobenius $(q$ -th) Π as an isogeny on Jacobian.

- Defn: • Isogeny $\alpha : J_C(\mathbb{K}) \rightarrow J_C(\mathbb{K})$ is a surjective morphism with finite $\ker(\alpha)$.
(respecting the group & the variety)

• deg $(\alpha) := |\ker(\alpha)|$.

- Ex. 1. $\Pi : J \rightarrow J$ is an isogeny with $\deg(\Pi) = 1$.

- Ex. 2. For $n \in \mathbb{Z}$, $[n] : J \rightarrow J ; D \mapsto n \cdot D$
is an isogeny. What's $\deg([n]) = ?$

- Defn: For prime $l \in \mathbb{N}$, we call $\ker([l])$ the l -torsion of \underline{J} , denoted $J[l]$.

• $T_l J$:= $\bigcup_{i \geq 1} J[l^i]$ is l -adic torsion of \underline{J} .

• Any isogeny $\alpha: J \rightarrow J$ gives an isogeny $T_l \alpha$: $T_l J \rightarrow T_l J$; $D \mapsto \alpha(D)$.

- Now, we can study $T_l \alpha$ for α coming from Π & $[n]$.

Theorem (Weil): (a) Linear map π on $T_e J$ has charpoly $t^{2g} L(1/t)$. (Assume $l \neq p$)

(b) $T_e J \cong (\mathbb{Z}_l)^{2g}$, i.e. $T_e J$ is a finite rank \mathbb{Z}_l -module.

(c) For any $n \in \mathbb{N}$, $\deg(\pi^n) = n^{2g}$ & $J(n) \cong (\mathbb{Z}/n)^{2g}$.

Pf: Recall, $\deg(\pi - 1) = |J(k)| = |\mathcal{C}_0(c)| = h(c)$
 $= L(1) = \prod_{i \in [2g]} (\alpha_i - 1)$.

• Similarly, $\forall m \in \mathbb{N}$, $\deg(\pi^m - 1) = |J(\mathbb{F}_q^m)|$
 $= \prod_{i=1}^{2g} (\alpha_i^m - 1)$ [Base-change of $L(t)$]

• Going to \mathbb{Z}_ℓ , we can view α_i in $\mathbb{C} \cap \overline{\mathbb{Q}}_\ell$.

• Let π act on $T_\ell J$ with eigenvalues $\beta_i, i \in \mathbb{N}$.

$$\Rightarrow \deg(\pi^m - 1) = |T_\ell J / (\pi^m - 1)T_\ell J| = \det(T_\ell(\pi^m - 1))$$

$$(|\ker| = |\text{coker}| = \det = \prod \text{eigenval}) = \prod_i (\beta_i^m - 1).$$

• Since, this holds $\forall m$, we deduce that

$$\alpha_i\text{'s} = \beta_i\text{'s} = \text{eigenvals of } T_\ell \pi \Rightarrow$$
$$\triangleright \text{charpoly}(\pi | T_\ell J) = t^{2g} \cdot L(1/t) \text{ over } \mathbb{Z}_\ell.$$

• Also, Δ charpoly $(\pi^{-1}|_{T_e J}) = q^{-g} \cdot L(t)$, over \mathbb{Z}_ℓ .

• (b) & (c) are implied by $\deg L = 2g$
& ℓ^i -torsion resp. n -torsion of J .

→ These properties are computationally useful, as one can try computing $L(t) \bmod \ell$ (or compute p -adically.) \square

– Current time-complexity for $L(t)$ is the min of:
 $\text{poly}(p, g, \delta)$ and $\text{poly}(\ell g p, 2^{2g}, \delta)$.
[kedlaya, 2001] [Pila, 1990]