## Program Analysis

https://www.cse.iitb.ac.in/~karkare/cs618/

## Welcome \& Introduction

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Program Analysis: About the course

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- Analysis of a Program, by a Program, for a Program
"Democracy is the government of the people, by the people, for the people" - Abraham Lincoln


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- by a Program - Analyzer (Compiler, Runtime)
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- of a Program - User Program
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- for a Program - Optimizer, Verifier, ...
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## Expectations from You

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-GCC, LLVM, SOOT


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- Understand and modify large code base -GCC, LLVM, SOOT
- Read state of the art research papers
-Discussions in class


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## -?

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- These can be announced or unannounced (surprize quizzes).
- Always bring a pen and some loose papers to the class


## QQ \#1 (Ungraded)



- What are the various phases of compilers that you know? (5 minutes)

Assignments / Exercises

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- 4-5 Assignments for the semester


## Using Program Analysis

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- Compiler Code Optimizations
- Why are optimizations important?
- Why not write optimized code to begin with?
- Where do optimizations fit in the compiler flow?


## Code Optimization

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- Machine Independent
- Remove redundancy introduced by the Programmer
- Remove redundancy not required by later phases of compiler
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Optimization must preserve the semantics of the original program!

## Machine Independent Optimizations

## Motivational Example

void quicksort(int m, int $n$ )
/* recursively sort a[m] through a[n] */
$\{$
int i, j;
int v, x;
if(n <= m) return;
i = m - 1; j = n; v = a[n];
while (1) \{
do i = i+1; while (a[i] < v);

$$
\text { do } j=j-1 ; ~ w h i l e ~(a[j]>v) ;
$$

if (i > j) break; $x=a[i] ; a[i]=a[j] ; a[j]=x ;$
\}
x = a[i]; $a[i]=a[n] ; a[n]=x ;$ quicksort(m,j); quicksort(i+1,n);
\}

## Motivational Example

void quicksort(int m, int n)
/* recursively sort a[m] through a[n] */
\{

```
int i, j;
int v, x;
if(n <= m) return;
\[
\begin{aligned}
& i=m-1 \cdot i=n \cdot v=a\lceil n\rceil . \\
& i=m-1 ; j=n ; v=a[n] ;
\end{aligned}
\]
while (1) \{
\[
\text { do } i=i+1 ; \text { while }(a[i]<v)
\]
\[
\text { do } j=j-1 ; \text { while }(a[j]>v) ;
\]
\[
\text { if }(i>j) \text { break; }
\]
\[
x=a[i] ; a[i]=a[j] ; a[j]=x ;
\]
\[
\}
\]
\[
x=a[i] ; a[i]=a[n] ; a[n]=x ;
\]
```

(1) $i=m-1$
(2) $\mathrm{j}=\mathrm{n}$
(3) $\mathrm{t} 1=4^{*} \mathrm{n}$
(4) $v=a[t 1]$
(5) i $=i+1$
(6) $\mathrm{t} 2=4^{*} \mathrm{i}$
(7) $\mathrm{t} 3=\mathrm{a}[\mathrm{t} 2]$
(8) if t3<v goto (5)
(9) $\mathrm{j}=\mathrm{j}-1$
(10) $\mathrm{t} 4=4^{*} \mathrm{j}$
(11) $\mathrm{t} 5=\mathrm{a}[\mathrm{t} 4]$
(12) if t5>v goto (9)
(13) if i>=j goto(23)
(14) t6 = 4*i
(15) $x=a[t 6]$
(16) t7 $=4^{* i}$
(17) $\mathrm{t} 8=4^{*} \mathrm{j}$
(18) t9 = a[t8]
(19) $\mathrm{a}[\mathrm{t} 7]=\mathrm{t} 9$
(20) t10 $=4^{*} j$
(21) $\mathrm{a}[\mathrm{t} 10]=\mathrm{x}$
(22) goto (5)
(23) t11 $=4 * i$
(24) $x=a[t 11]$
(25) t12 $=4 * i$
(26) t13 $=4^{*}$ n
(27) $\mathrm{t} 14=\mathrm{a}[\mathrm{t} 13]$
(28) $\mathrm{a}[\mathrm{t} 12]=\mathrm{t} 14$
(29) t15 $=4 * n$
(30) $a[t 15]=x$

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(15) $x=a[t 6]$
(16) t7 $=4^{* i}$
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(21) $\mathrm{a}[\mathrm{t} 10]=\mathrm{x}$
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(24) $x=a[t 11]$
(25) t12 $=4 * i$
(26) t13 $=4^{*}$ n
(27) t14 = a[t13]
(28) $\mathrm{a}[\mathrm{t} 12]=\mathrm{t} 14$
(29) t15 $=4 * n$
(30) $a[t 15]=x$

(14) t6 = 4*i
(15) $\mathrm{x}=\mathrm{a}[\mathrm{t} 6]$
(16) $\mathrm{t} 7=4^{*} \mathrm{i}$
(17) $\mathrm{t} 8=4^{*} \mathrm{j}$
(18) t9 = a[t8]
(19) $\mathrm{a}[\mathrm{t} 7]=\mathrm{t} 9$
(20) t10 $=4^{*} j$
(21) $\mathrm{a}[\mathrm{t} 10]=\mathrm{x}$
(22) goto (5)
(23) t11 $=4 * i$
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(30) $\mathrm{a}[\mathrm{t} 15]=\mathrm{x}$

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(22) goto (5)
(23) t11 $=4^{*}$ i
(24) $x=a[t 11]$
(25) t12 = 4*i
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(17) t8 $=4^{*} \mathrm{j}$
(18) $\mathrm{t} 9 \mathrm{=} \mathrm{a}[\mathrm{t} 8]$
(19) $\mathrm{a}[\mathrm{t} 7]=\mathrm{t} 9$
(20) t10 $=4^{*} j$
(21) $\mathrm{a}[\mathrm{t} 10]=\mathrm{x}$
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| (1) | $i=m-1$ |
| :---: | :---: |
| (2) | $j=n$ |
| (3) | t1 $=4^{*} n$ |
| (4) | $\mathrm{v}=\mathrm{a}[\mathrm{t} 1]$ |
| (5) | $i=1+1$ |
| (6) | t2 $=4^{*}$ i |
| (7) | $\mathrm{t} 3=\mathrm{a}$ [t2] |
| (8) | if $t 3<v$ goto (5) |
| (9) | $\mathrm{J}=\mathrm{J}-1$ |
| (10) | t4 $=4^{*} j$ |
| (11) | $t 5=a[t 4]$ |
| (12) | if $\mathrm{t} 5>\mathrm{V}$ goto (9) |
| (13) | if i>=j goto(23) |


| (14) | $t 6=4^{*} i$ |
| :---: | :---: |
| (15) | $x=a[t 6]$ |
| (16) | $t 7=4^{*} i$ |
| (17) | $t 8=4^{*} j$ |
| (18) | $t 9=a[t 8]$ |
| (19) | $a[t 7]=t 9$ |
| (20) | $t 10=4^{*} j$ |
| (21) | $a[t 10]=x$ |
| (22) | goto (5) |
| (23) | t11 $=4 * i$ |
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| (25) | $t 12=4 * i$ |
| (26) | $t 13=4^{*} n$ |
| (27) | $t 14=a[t 13]$ |
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## Common Subexpression Elimination



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Common Subexpression Elimination


## Did I miss one common subexpression?



Elimination not safe as $\mathbf{a}$ [] is modified on path

$$
\mathrm{B} 1->\mathrm{B} 2->\mathrm{B} 3->\mathrm{B} 4->\mathrm{B} 5->\mathrm{B} 2->\mathrm{B} 3->\mathrm{B} 4->\mathrm{B} 6
$$



## Copy Propagation



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## Dead Code Elimination



Strength Reduction


Strength Reduction


## Induction Variable elimination



## Induction Variable elimination



Dead Code elimination (again!)


## Benefits

| Block | \# Stmt before <br> Optimizations | \# Stmt after <br> Optimizations |
| :--- | :--- | :--- |
| B1 | 4 | 6 |
| B2 | 4 | 3 |
| B3 | 4 | 3 |
| B4 | 1 | 1 |
| B5 | 9 | 3 |
| B6 | 8 | 3 |

## Benefits

| Block | \#Stmt before <br> Optimizations | \# Stmt after <br> Optimizations |
| :--- | :--- | :--- |
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| B2 | 4 | 3 |
| B3 | 4 | 3 |
| B4 | 1 | 1 |
| B5 | 9 | 3 |
| B6 | 8 | 3 |

Assume: Unit cost for statements, 10 iterations of outer loop, 100 iterations of each inner loop
Cost of execution:
ORIGINAL: $1 * 4+100 * 4+100 * 4+10 * 1+10 * 9+1 * 8=912$
OPTIMIZED: $1 * 6+100 * 3+100 * 3+10 * 1+10 * 3+1 * 3=649$

# Machine Dependent Optimizations 

## Peephole Optimization

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- target code often contains redundant instructions and suboptimal constructs


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- target code often contains redundant instructions and suboptimal constructs
- examine a short sequence of target instruction (peephole) and replace by a shorter or faster sequence
- peephole is a small moving window on the target systems


# Peephole optimization examples... 

## Redundant loads and stores

- Consider the code sequence

Move $\mathrm{R}_{0}$, a
Move a, $\mathrm{R}_{0}$

# Peephole optimization examples... 

Redundant loads and stores

- Consider the code sequence

Move $\mathrm{R}_{0}$, a
Move a, $\mathrm{R}_{0}$

- Instruction 2 can always be removed if it does not have a label.


## Peephole optimization examples...

## Unreachable code

- Consider following code sequence \#define debug 0
if (debug) \{
print debugging info
\}


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- Consider following code sequence \#define debug 0
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this may be translated as
if debug $==1$ goto L1
goto L2
L1: print debugging info
L2:


## Peephole optimization examples...

## Unreachable code

- Consider following code sequence
\#define debug 0
if (debug) \{
print debugging info
\}
this may be translated as
if debug == 1 goto L1
goto L2
L1: print debugging info
L2:

Eliminate jumps
if debug != 1 goto L2
print debugging information

## Unreachable code example ...

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constant propagation

$$
\text { if } 0 \text { <> } 1 \text { goto L2 }
$$

print debugging information
L2:

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$$
\text { if } 0 \text { <> } 1 \text { goto L2 }
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print debugging information
L2:

Evaluate boolean expression. Since if condition is always true the code becomes
goto L2
print debugging information
L2:

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constant propagation

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\text { if } 0 \text { <> } 1 \text { goto L2 }
$$

print debugging information
L2:

Evaluate boolean expression. Since if condition is always true the code becomes
goto L2
print debugging information
L2:

The print statement is now unreachable. Therefore, the code becomes

L2:

## Peephole optimization examples...

- flow of control: replace jump over jumps


L1 : goto L2

## Peephole optimization examples...

- flow of control: replace jump over jumps
goto L1
goto L2
by

L1: goto L2
L1 : goto L2

- Simplify algebraic expressions remove $x:=x+0$ or $x:=x * 1$


## Peephole optimization examples...

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- Strength reduction
- Replace $\mathrm{X}^{\wedge} 2$ by $\mathrm{X}^{*} \mathrm{X}$
- Replace multiplication by left shift
- Replace division by right shift


## Peephole optimization examples...

- Strength reduction
- Replace $\mathrm{X}^{\wedge} 2$ by $\mathrm{X}^{*} \mathrm{X}$
- Replace multiplication by left shift
- Replace division by right shift
- Use faster machine instructions

$$
\begin{array}{ll}
\text { replace } & \text { Add \#1,R } \\
\text { by } & \text { Inc R }
\end{array}
$$

## Course Logistics

## Proposed Evaluation

| Assignments | 5\%-10\% |
| :--- | ---: |
| Course Project | 30\%-40\% |
| ( Proposal | 5\% ) |
| ( Report | 15\% ) |
| ( Presentation | $15 \%$ ) |
| Mid semester exam | $\mathbf{1 0 \% - 2 0 \%}$ |
| End semester exam | $\mathbf{2 5 \% - 3 5 \%}$ |
| Quizzes/Class Participation | $\mathbf{5 \%}$ |

Project

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- You need to implement some non-trivial analysis/optimization using one of the open source infrastructure
- For e.g., some paper published in last 10 years
- You are encouraged to suggest your own projects
- Bonus marks for publishable results
- Individual OR Group of 2

Assignment \#1

## Assignment \#1

- Select one of the compiler infrastructure mentioned on the course webpage and
a) Download it
b) Build it
c) Submit a report
d) one page about the infrastructure, and the optimizations present in it.
e) one page about the most interesting optimization found, with example


## Assignment \#1

- You can try more than one tool, even something not mentioned on the webpage.
- But submit report for only one.
- Preferably the one you plan to use for your project.


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- DEADLINE: July 30th, End of Day (before Midnight)
- See course website for submission details (TBD)
https://www.cse.iitb.ac.in/~karkare/cs618/

