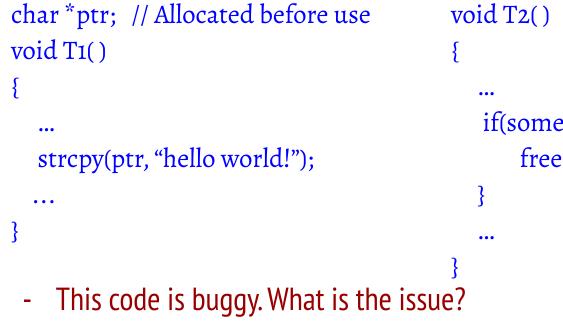
CS330: Operating Systems

Concurrency bugs

Common issues in concurrent programs

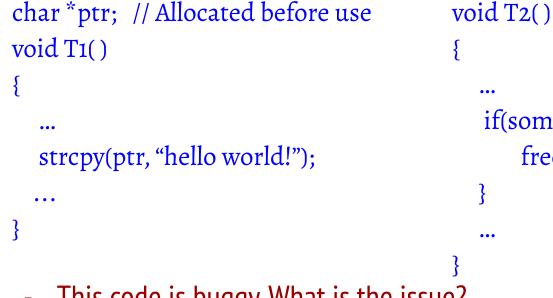
- Atomicity issues
- Failure of ordering assumption
- Deadlocks

Concurrency bugs - atomicity issues



... if(some_condition){ free(ptr); ptr = NULL;

Concurrency bugs - atomicity issues

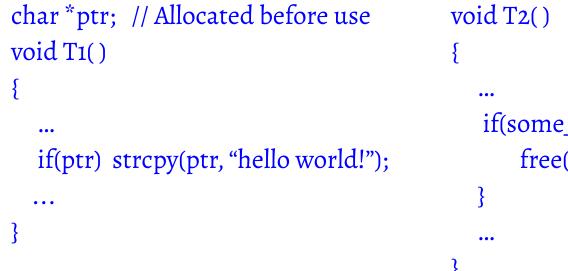


... if(some_condition){ free(ptr); ptr = NULL;

- This code is buggy. What is the issue?
- T2 can free the pointer before T1 uses it.
- How to fix it?

```
Concurrency bugs - atomicity issues
                                            void T<sub>2</sub>()
char *ptr; // Allocated before use
void T<sub>1</sub>()
                                               ...
                                               if(some_condition){
   ...
                                                   free(ptr); ptr = NULL;
  if(ptr) strcpy(ptr, "hello world!");
  . . .
                                               ...
    Does the above fix (checking ptr in T1) work?
 -
```

Concurrency bugs - atomicity issues



if(some_condition){ free(ptr); ptr = NULL;

- Does the above fix (checking ptr in T1) work?
- Not really. Consider the following order of execution:
- T1: "if(ptr)" T2: "free(ptr)" T1: "strcpy" Result: Segfault

Concurrency bugs - ordering issues

- 1. bool pending;
- 2. void T1()
- 3. {

7. }

- 4. pending = true;
- 5. do_large_processing();
- 6. while (pending);

- 1. void T2()
- 2. {
- 3. do_some_processing();
- 4. pending = false;
- 5. some_other_processing();

6. }

- This code works with the assumption that line#4 of T2 is executed after line#4 of T1
- If this ordering is violated, T1 is stuck in the while loop

Concurrency bugs - deadlocks

```
struct acc t{
     lock_t *L;
     id_t acc_no;
     long balance;
}
void txn_transfer( acc_t *src,
                 acc_t *dst, long amount)
  lock(src \rightarrow L); lock(dst \rightarrow L);
  check_and_transfer(src, dst, amount);
  unlock(dst \rightarrow L); unlock(src \rightarrow L);
```

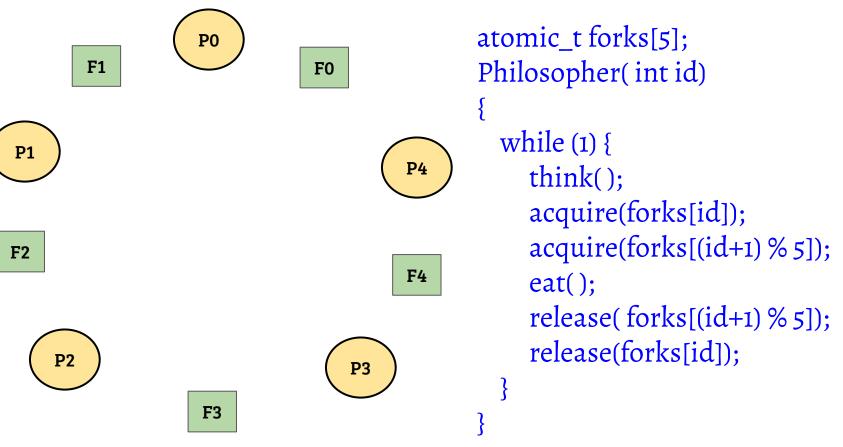
- Consider a simple transfer transaction in a bank
- Where is the deadlock?

Concurrency bugs - deadlocks

```
struct acc_t{
     lock_t *L;
     id_t acc_no;
     long balance;
}
void txn_transfer( acc_t *src,
                acc_t *dst, long amount)
  lock(src \rightarrow L); lock(dst \rightarrow L);
  check_and_transfer(src, dst, amount);
  unlock(dst \rightarrow L); unlock(src \rightarrow L);
```

- Consider a simple transfer transaction in a bank
- Where is the deadlock?
- T1: txn_transfer(iitk, cse, 10000)
 - lock (iitk), lock (cse)
- T2: txn_transfer(cse, iitk, 5000)
 - lock (cse), lock(iitk)

Dining philosophers



Conditions for deadlock

- Mutual exclusion: exclusive control of resources (e.g, thread holding lock)
- Hold-and-wait: hold one resource and wait for other
- No resource preemption: Resources can not be forcibly removed from threads holding them
- Circular wait: A cycle of threads requesting locks held by others. Specifically, a cycle in the directed graph G (V, E) where V is the set of processes and $(v1, v2) \in E$ if v1 is waiting for a lock held by v2

All of the above conditions should be satisfied for a deadlock to occur

Solutions for deadlocks

- Remove mutual exclusion: lock free data structures
- Either acquire all resources or no resource
 - trylock(lock) APIs can be used (e.g., pthread_mutex_trylock())
- Careful scheduling: Avoid scheduling threads such that no deadlock occur
- Most commonly used technique is to avoid circular wait. This can be achieved by ordering the resources and acquiring them in a particular order from all the threads.

Concurrency bugs - avoiding deadlocks

```
struct acc_t{
```

```
lock_t *L;
     id_t acc_no;
     long balance;
}
void txn_transfer( acc_t *src,
                acc_t *dst, long amount)
  lock(src \rightarrow L); lock(dst \rightarrow L);
```

```
check_and_transfer(src, amount);
unlock(dst → L); unlock(src → L);
```

- Deadlock in a simple transfer transaction in a bank
- While acquiring locks, first acquire the lock for the account with lower "acc_no" value
- Account number comparison
 performed before acquiring the lock

Dining philosophers: breaking the deadlock

